# **Dynamical System Decomposition Using Dissipation Inequalities**

James Anderson

André Teixeira

Henrik Sandberg

Antonis Papachristodoulou

Abstract— In this paper we investigate stability and interaction measures for interconnected systems that have been produced by decomposing a large-scale linear system into a set of lower order subsystems connected in feedback. We begin by analyzing the requirements for asymptotic stability through generalized dissipation inequalities and storage functions. Using this insight we then describe various metrics based on a system's energy dissipation to determine how strongly the subsystems interact with each other. From these metrics a decomposition algorithm is described.

#### I. INTRODUCTION

The scalable analysis and design of complex dynamical systems is a challenging area in systems and control theory. In this work we describe a set of algorithms that can be used to analyze the stability and characterize the interconnection strength of Linear Time Invariant (LTI) dynamical systems. The methods proposed are based on the notion of dynamical system decomposition [1] and dissipation inequalities with quadratic supply rates [2].

It is frequently the case that many systems have an underlying network structure. If the network structure or connection topology is known *a priori* then this information can be used to help design scalable algorithms for interrogating the system of interest. When the network structure is not known it is important to impose an interconnection topology (decomposition) in order to facilitate further analysis.

In this paper two issues are addressed. We begin by deriving stability criteria for interconnected LTI subsystems using dissipation inequalities and quadratic supply rate functions [2], [3]. The subsystems of interest may have been obtained through a decomposition algorithm or they may have a physical realization. In the sequel a method for decomposing networks using the supply rate as a metric for interconnection strength is described and illustrated on an RC network.

System decomposition was first suggested as a framework for handling large-scale systems by Šiljak [4]. However the framework did not provide any insight on how to produce the system decomposition. In recent work [1], [5] an algorithmic method for producing decompositions based on representing the system as a graph and minimizing the worst case "energy flow" between states was presented. In [6] an alternative approach using Hankel-norm based lumping technique for decomposition was presented.

J.A and A.T contributed equally to this work. J. Anderson and A. Papachristodoulou are with the Department of Engineering Science, University of Oxford, Parks Road, Oxford, OX1 3PJ. {james.anderson,antonis}@eng.ox.ac.uk

A. Teixeira and H. Sandberg are with the Automatic Control Laboratory, School of Electrical Engineering, KTH-Royal Institute of Technology, Osquldas vg 10, SE-10044 Stockholm, Sweden. {andretei,hsan}@kth.se Using the framework in [1], to decompose a dynamical system model, the system must first be represented as a graph where each node is a state of the system and weighted edges represent state interactions. The decomposition is realized by applying a spectral graph partitioning algorithm [7], [8] to the weighted graph. Here we extend this idea to the case where nodes may represent subsystems and edges represent the strength of interaction between subsystems as determined by a given supply rate function. We investigate the physical interpretation of a class of supply rates and conclude by presenting a clustering based decomposition algorithm.

The paper is organized as follows: In Section II we introduce the necessary background material. Section III uses quadratic supply rate functions to derive stability criteria using passivity, input strict passivity and finite-gain arguments. In Section IV various edge weight metrics are described and a decomposition algorithm presented. A numerical example is given in Section V and the paper is concluded in Section VI.

#### **II. PRELIMINARIES**

#### A. Notation

 $\mathbb{R}^n$  denotes the *n*-dimensional Euclidean space,  $\mathbb{R}^{n \times m}$  denotes the set of  $n \times m$  real matrices. If  $M \in \mathbb{R}^{n \times n}$  and  $M = M^{\top}$  then M > 0,  $M \ge 0$  denote that M is positive definite, positive semidefinite respectively. The maximum singular value of M is denoted by  $\bar{\sigma}(M)$ . Given k matrices  $M_1, \ldots, M_k$ , diag $(M_1, \ldots, M_k)$  denotes the concatenated block diagonal matrix.

#### B. Storage Functions and Quadratic Supply

We consider dissipation inequalities containing quadratic nonnegative storage functions, quadratic supply rates, and the notion of (Q, S, R) dissipativity described in [3]. Consider the LTI system

$$\frac{d}{dt}x(t) \triangleq \dot{x}(t) = Ax(t) + Bu(t), \quad y(t) = Cx(t), \quad (1)$$

where  $x(t) \in \mathbb{R}^n$ ,  $u(t) \in \mathbb{R}^p$ , and  $y(t) \in \mathbb{R}^m$  are the state, input and output vectors, respectively. For simplicity we will omit the time argument and refer to the former as simply x, u, and y.

System (1) is said to be *dissipative* with respect to the supply rate w(u, y) if there exists a continuously differentiable *storage* function  $V : \mathbb{R}^n \to \mathbb{R}$  satisfying the following *dissipation inequality* 

$$\dot{V}(x) \le w(u, y),\tag{2}$$

with V(0) = 0 and  $V(x) \ge 0$  for all  $x \ne 0$ . We are interested in quadratic supply rate functions of the form

$$w(u,y) = y^{\top}Qy + 2u^{\top}Sy + u^{\top}Ru, \qquad (3)$$

with  $Q = Q^{\top}$  and  $R = R^{\top}$ , where u and y form input output pairs (u, y) and Q, S and R are of appropriate dimension. The LTI system (1) is said to be dissipative if the following holds

$$\begin{split} V(x) &= x^\top P x > 0 \quad \text{for all } x \neq 0 \\ \dot{V}(x) &= x^\top (A^\top P + P A) x + u^\top B^\top P x + x^\top P B u \\ &\leq w(u,y) \end{split}$$

where  $P = P^{\top}$ . Dissipativity can be checked by solving the following LMI

$$\begin{bmatrix} A^{\top}P + PA - C^{\top}QC & PB - C^{\top}S \\ B^{\top}P - S^{\top}C & -R \end{bmatrix} \le 0, \quad (4)$$

for P > 0.

# C. Stability of Interconnected Dissipative Systems

We study the interconnection of N subsystems,  $\Sigma_i$ , where

$$\Sigma_i \begin{cases} \dot{x}_i = A_i x_i + B_i u_i \\ y_i = C_i x_i, \end{cases}$$
(5)

and the interconnection between subsystems is

$$u_i = -\sum_{j=1}^N H_{ij} y_j.$$
 (6)

Defining  $x = [x_1^\top \cdots x_N^\top]^\top$ ,  $u = [u_1^\top \cdots u_N^\top]^\top$ , and  $y = [y_1^\top \cdots y_N^\top]^\top$ , the interconnection may be written as u = -Hy and the global system dynamics are described by

$$\dot{x} = Ax - BHCx, \quad y = Cx, \tag{7}$$

where  $A = \text{diag}(A_1, \dots, A_N)$ ,  $B = \text{diag}(B_1, \dots, B_N)$ , and  $C = \text{diag}(C_1, \dots, C_N)$ .

Assumption 1: Each subsystem  $\Sigma_i$  is dissipative with respect to a given supply rate  $(Q_i, S_i, R_i)$ .

Given the previous assumption, for each subsystem there exists a symmetric matrix  $P_i > 0$  such that the dissipation inequality (2) holds with the  $(Q_i, S_i, R_i)$  supply rate. Therefore the inequality

$$\sum_{i=1}^{N} \dot{V}_i(x_i) \le \sum_{i=1}^{N} w_i(u_i, y_i)$$
(8)

holds regardless of the interconnection.

Remark 1: Defining  $P = \operatorname{diag}(P_1, \dots, P_N)$  and  $V(x) = x^\top P x$ , we have  $\dot{V}(x) = x^\top (A^\top P + PA)x + u^\top B^\top P x + x^\top P B u = \sum_{i=1}^N \dot{V}_i(x_i)$ . Furthermore,  $w(u, y) = \sum_{i=1}^N w_i(u_i, y_i)$  is given by (3) with  $Q = \operatorname{diag}(Q_1, \dots, Q_N)$ ,  $S = \operatorname{diag}(S_1, \dots, S_N)$ , and  $R = \operatorname{diag}(R_1, \dots, R_N)$ .

Inequality (8) may then be rewritten as

$$\sum_{i=1}^{N} \dot{V}_i(x_i) \le y^{\top} \hat{Q} y, \tag{9}$$

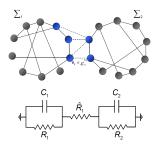


Fig. 1. Top: Graph of interconnected subsystems  $\Sigma_1, \Sigma_2$ . Boundary nodes  $\mathcal{V}_b^{12}$  are blue, dashed edges,  $\mathcal{E}_c^{12}$ , connect the subsystems via the boundary nodes. Bottom: Two nodes in an RC network, edge  $e_k \in \mathcal{E}_c^{12}$  corresponds to resistor  $\widehat{R}_1$  and the vertices correspond to the resistor and capacitors connected in parallel.

where  $\hat{Q} = Q - S^{\top}H - H^{\top}S + H^{\top}RH$ . A sufficient condition for stability of the global interconnected system follows from (9):

Lemma 1 ([3]): Assume each system  $\Sigma_i$  is observable. The global interconnected system is asymptotically stable if  $\hat{Q}$  is negative definite. Furthermore,  $V(x) = \sum_{i=1}^{N} V_i(x_i) = x^{\top} Px$  is a Lyapunov function for the global system, with  $P = \text{diag}(P_1, \dots, P_N)$ .

# D. Algebraic Graph Theory

Consider an *undirected weighted* graph  $\mathcal{G}(\mathcal{V}, \mathcal{E}, \mathcal{Z})$  where  $\mathcal{V} = \{v_1, \dots, v_N\}$  is the set of N vertices or nodes,  $\mathcal{E} = \{e_1, \dots, e_M\} \subseteq \mathcal{V} \times \mathcal{V}$  is the edge-set and  $\mathcal{Z} = \{z_1, \dots, z_M\}$  where  $z_j > 0$  is the weight of edge j. Associated with  $\mathcal{G}$  is a symmetric weighted adjacency matrix  $\mathcal{A}(\mathcal{G}) \in \mathbb{R}^{N \times N}$  where  $[\mathcal{A}(\mathcal{G})]_{ij} > 0$  if there exists an edge connecting  $v_i$  to  $v_j$ . The weighted  $N \times N$  Laplacian matrix is defined by  $\mathcal{L}(\mathcal{G}) = \text{diag}(\mathcal{A}(\mathcal{G})\mathbf{1}) - \mathcal{A}(\mathcal{G})$ . For undirected graphs  $\mathcal{L}(\mathcal{G}) \geq 0$  and the number of zero eigenvalues is equal to the number of connected components in the graph.

The incidence matrix,  $C(\mathcal{G}) \in \mathbb{R}^{N \times M}$  of an undirected graph is defined by assigning an arbitrary direction to each  $e_i \in \mathcal{E}$  and setting  $[C(\mathcal{G})]_{ij} =$ 1 if  $e_i$  enters  $v_j$ , -1 if  $e_i$  leaves  $v_j$  and 0 otherwise. The weighted Laplacian can then be equivalently defined by  $\mathcal{L}(\mathcal{G}) = C(\mathcal{G})W(\mathcal{G})C(\mathcal{G})^{\top}$  where  $W(\mathcal{G}) \in \mathbb{R}^{M \times M}$  is a diagonal matrix and  $[W(\mathcal{G})]_{ii} = z_i$  with  $z_i \in \mathcal{Z}$ . Given a graph  $\mathcal{G}(\mathcal{V}, \mathcal{E}, \mathcal{Z})$ , consider a subset of the vertices  $\mathcal{V}_j \subset \mathcal{V}$ we call the graph  $\mathcal{G}_j = \mathcal{G}/\mathcal{V}_j$  an *induced subgraph* of  $\mathcal{G}$ .

Assume that we have a graph  $\mathcal{G}(\mathcal{V}, \mathcal{E}, \mathcal{Z})$  which has been partitioned such that  $\mathcal{V} = \mathcal{V}_1 \bigcup \mathcal{V}_2$  and  $\mathcal{V}_1 \bigcap \mathcal{V}_2 = \emptyset$ . If there exists an edge  $e_k = (v_i, v_j) \in \mathcal{E}$  such that  $v_i \in \mathcal{V}_1$  and  $v_i \in \mathcal{V}_2$  (or vice-versa) then  $v_i$  and  $v_j$  are called boundary nodes and belong to the set  $\mathcal{V}_b^{12}$ . The set of edges that connect boundary nodes in  $\mathcal{V}_1, \mathcal{V}_2$  is given by  $\mathcal{E}_c^{12}$ .

When it is clear from the context we will omit the graph argument ands refer to the adjacency, incidence, weighting and Laplacian matrices by  $\mathcal{A}, \mathcal{C}, \mathcal{W}$  and  $\mathcal{L}$  respectively. For a thorough overview of algebraic graph theory see [9].

#### E. Illustrative Example

The ideas presented above are now illustrated through a resistor capacitor network example. The network is represented by the undirected weighted graph  $\mathcal{G}(\mathcal{V}, \mathcal{E}, \mathcal{Z})$ . In this setting, each vertex  $v_i \in \mathcal{V}$  corresponds to a capacitor  $C_i$  in parallel with a resistor  $R_i$  connected to ground and each edge  $e_k \in \mathcal{E}$ represents a resistor  $\hat{R}_k$  with  $z_k = 1/\hat{R}_k \in \mathcal{Z}$  connecting the free terminals of two vertices. This is illustrated in Figure 1. Let all capacitors have unit capacitance and denote  $x_i \in \mathbb{R}$ as the voltage in capacitor i and  $u_i \in \mathbb{R}$  the current entering node i. Each node is modeled as by the first-order system

$$\dot{x}_i = -g_i x_i + u_i, \quad y_i = x_i,$$
 (10)

where  $g_i = 1/R_i$ . Defining  $G = \text{diag}(g_1, \dots, g_N)$  the dynamics of the global network are given by

$$\dot{x} = -(\mathcal{L} + G)x, \quad y = x. \tag{11}$$

Let  $\Sigma_j$  be a subnetwork described by the induced graph  $\mathcal{G}_j(\mathcal{V}_j, \mathcal{E}_j, \mathcal{Z}_j)$ . The dynamics of  $\Sigma_j$  are described by

$$\Sigma_j \begin{cases} \dot{x}_j = -(\mathcal{L}_j + G_j)x_j + B_j u_j \\ y_j = B_j^\top x_j \end{cases}$$
(12)

where  $u_j$  represents the input current to  $\Sigma_j$  from the rest of the network that enters through the boundary nodes whose dynamics are described by the matrix  $B_j$ . To construct  $B_j$ : i) Define the following sets  $\mathcal{V}_i, \mathcal{V}_j, \mathcal{E}_i, \mathcal{E}_j, \mathcal{V}_b^{ij}$  and  $\mathcal{E}_c^{ij}$ . ii) From these sets construct an incidence matrix corresponding to  $\mathcal{V}_b^{ij}, \mathcal{E}_c^{ij}$  such that all edges enter  $\mathcal{V}_j$ , iii)  $B_j$  corresponds to the part of the incidence matrix with nodes belonging to  $\mathcal{V}_j$ .

Defining the storage function for this system as  $V_j(x_j) = \frac{1}{2}x_j^{\top}x_j$  we have  $\dot{V}_j(x_j) = -x_j^{\top}\mathcal{L}_jx_j - x_j^{\top}G_jx_j + u_j^{\top}y_j$ , where the power dissipated on the internal edge resistors, the power dissipated on the node resistors, and the input power to  $\Sigma_j$  correspond to  $-x_j^{\top}\mathcal{L}_jx_j, -x_j^{\top}G_jx_j$ , and  $u_j^{\top}y_j$ , respectively.

Given  $V_j(x_j)$ , two interesting supply rate functions for which the system is dissipative can be immediately identified.

Observation 1: If  $\Sigma_j$  is stable, then  $\Sigma_j$  is (0, I, 0)dissipative. Furthermore, consider  $x_j^{\top}G_jx_j = x_j^{\top}\tilde{G}_jx_j + y_j^{\top}G_j^by_j$ , where the first term is the power dissipated in the internal node

the first term is the power dissipated in the internal node resistors, while the second term corresponds to power dissipated on the boundary node resistors.

Observation 2: If  $\mathcal{L}_i + \tilde{G}_j$  is positive semidefinite, then  $\Sigma_j$  is  $(-G_j^b, I, 0)$ -dissipative.

#### **III. STABILITY ANALYSIS**

The goal of system decomposition is to provide tractable means of analyzing dynamical systems that have a large state dimension. The idea is to take an LTI system model with high state dimension and *decompose* it into the form of (5) where composite methods can then be used to infer stability of the original system.

Here we describe some stability results based on composite Storage functions. We assume that a decomposition algorithm has already been applied to obtain the subsystems. In Section IV we present a clustering algorithm for decomposition. Consider the LTI system  $\Sigma$ :

$$\Sigma \begin{cases} \dot{x} = Ax, \quad x(0) = x_0, \\ y = x \end{cases}$$
(13)

with  $x \in \mathbb{R}^n$ . Now assume that it has been decomposed into two subsystems connected in feedback

$$\Sigma_{1} \begin{cases} \dot{x}_{1} = A_{11}x_{1} + u_{1} \\ u_{1} = A_{12}x_{2} \\ y_{1} = x_{1} \end{cases}, \\ \Sigma_{2} \begin{cases} \dot{x}_{2} = A_{22}x_{2} + u_{2} \\ u_{2} = A_{21}x_{1} \\ y_{2} = x_{2} \end{cases}$$
(14)

where the state vector has been permuted such that  $x = [x_1^{\top}, x_2^{\top}]^{\top}$  and  $x_1 \in \mathbb{R}^{n_1}, x_2 \in \mathbb{R}^{n_2}, n_1 + n_2 = n$  and no state belongs to multiple subsystems.

*Remark 2:* Assume system  $\Sigma$  has been decomposed into  $\Sigma_1$  and  $\Sigma_2$  which are dissipative w.r.t. the quadratic supply rates  $w_1(u_1, y_1)$  and  $w_2(u_2, y_2)$  of the form (3) respectively. If  $w_1(u_1, y_1) + w_2(u_2, y_2) < 0$  for all input output pairs then the sum of the Storage functions  $V_1(x_1) + V_2(x_2)$  is a Lyapunov function that proves that the equilibrium point of (13) is asymptotically stable. This is a direct application of Lemma 1.

Note that this and all further results generalize to the case where  $\Sigma$  is decomposed into multiple subsystems, for the sake of clarity we focus on the case of two subsystems.

If we assume a generic interconnection structure for  $\Sigma_1, \Sigma_2$  of the form

$$\begin{bmatrix} u_1 \\ u_2 \end{bmatrix} = \begin{bmatrix} H_{11} & H_{12} \\ H_{21} & H_{22} \end{bmatrix} \begin{bmatrix} y_1 \\ y_2 \end{bmatrix}$$
(15)

then the right hand side of  $\dot{V}_1(x_1) + \dot{V}_2(x_2) \le w_1(u_1, y_1) + w_2(u_2, y_2)$  with  $(u_i, y_i)$  obtained from the decomposition and interconnection matrix (15) is given by (16) on the next page. By appropriate choice of the Q, S, R matrices, the supply functions (3) can represent passivity, finite-gain and dissipativity, each of which alters the structure of (16). The remainder of this section examines each of these cases in turn and provides stability tests for (13) based on its decomposed subsystems (14).

#### A. Passivity

An LTI system of the form (5) is said to be passive if it is dissipative with respect to supply rate (3) with  $(Q_i, S_i, R_i) = (0, I, 0)$  and LMI (4) is feasible. Assume that  $\Sigma$  has been decomposed into  $\Sigma_1, \Sigma_2$  (which is equivalent to (5)). Substituting the appropriate matrices into the supply rate functions, we see from Equation (16) that we require

$$\begin{bmatrix} H_{11}^{\top}S_1 + S_1^{\top}H_{11} & S_1^{\top}H_{12} + H_{21}^{\top}S_2 \\ \star & H_{22}^{\top}S_2 + S_2^{\top}H_{22} \end{bmatrix} < 0, \quad (17)$$

where from (14) we have that

$$H = \left[ \begin{array}{cc} 0 & A_{12} \\ A_{21} & 0 \end{array} \right].$$

With this interconnection structure and system decomposition the diagonal block entries in (17) are zero and the off diagonal blocks are given by  $A_{12} + A_{21}^{\top}$  and its transposition. In this form (17) cannot be negative definite as its eigenvalues

$$\begin{bmatrix} x_{1} \\ x_{2} \end{bmatrix}^{\top} \begin{bmatrix} (H_{11}^{\top}R_{1}H_{11} + H_{11}^{\top}S_{1} + S_{1}^{\top}H_{11} + (H_{11}^{\top}R_{1}H_{12} + S_{1}^{\top}H_{12} + H_{12} + H_{12}^{\top}R_{2}H_{21} + Q_{1}) \\ + H_{21}^{\top}R_{2}H_{21} + Q_{1}) \\ \star (H_{22}^{\top}R_{2}H_{22} + H_{22}^{\top}S_{2} + S_{2}^{\top}H_{22} + H_{22}^{\top}H_{22} + H_{12}^{\top}R_{1}H_{12} + Q_{2}) \end{bmatrix} \begin{bmatrix} x_{1} \\ x_{2} \end{bmatrix}$$
(16)

will be real and symmetric about the imaginary axis. This problem can be alleviated if we consider a slight modification to the decomposition described by (14) by imposing a further decomposition on the drift matrices  $A_{ii}$  and include a feedback term. The new decomposition for  $\Sigma_1$  is

$$\widehat{\Sigma}_{1} \begin{cases} \dot{x}_{1} = \epsilon_{1}A_{11}x_{1} + u_{1} \\ u_{1} = A_{12}x_{2} + \delta_{1}A_{11}x_{1} \\ y_{1} = x_{1} \end{cases}$$
(18)

where  $\epsilon_1 + \delta_1 = 1$  and we assume all matrices are of compatible dimension. In the same manner  $\hat{\Sigma}_2$  can be constructed. LMI (17) is then replaced by

$$\begin{bmatrix} \delta_1 (A_{11} + A_{11}^{\top}) & A_{12} + A_{21}^{\top} \\ \star & \delta_2 (A_{22} + A_{22}^{\top}) \end{bmatrix} < 0$$
(19)

When  $\hat{\Sigma}_1, \hat{\Sigma}_2$  are dissipative with respect to (0, I, 0) and LMI (19) is feasible the original system (13) is stable as verified by the Lyapunov function  $V(x) = V_1(x_1) + V_2(x_2)$ . An alternative approach is to select  $\epsilon_i, \delta_i$  arbitrarily (ensuring  $\epsilon_1 + \delta_1 = 1$ ) and using the modified decomposition  $\hat{\Sigma}$  solve LMI (17) where the decision variables are the diagonal matrices  $S_i > 0$ . Such an approach is possible because any system that is dissipative with respect to (0, I, 0) is also dissipative with respect to any (0, X, 0) supply rate with X > 0 diagonal.

# B. Finite Gain

For LTI systems the  $\mathcal{L}_2$  gain from input to output of a system in the form of (5) can be calculated by solving:

$$\begin{array}{ll} \min & \gamma_i \\ \text{s.t.} & \begin{bmatrix} A_i^\top P_i + P_i A_i + C_i^\top C_i & P_i B_i \\ B_i^\top P_i & -\gamma_i^2 I \end{bmatrix} \leq 0 \ (20) \\ P_i > 0, \quad \gamma_i > 0. \end{array}$$

The  $\mathcal{L}_2 \to \mathcal{L}_2$  gain is then given by  $\gamma_i$  [10]. For two systems connected in feedback if  $\gamma_1\gamma_2 < 1$  then the feedback connection is stable [11]. A generalization of the small gain theorem for networks is given in [12], [13]. Following from LMI (20) it can be seen that the supply rate functions associated with finite gain are given by  $(-I, 0, \gamma_i^2 I)$  for i = 1, 2.

Substituting the appropriate Q, S, R matrices and interconnections into (16) gives the following stability requirement:

$$\begin{bmatrix} \gamma_2^2 A_{21}^{\dagger} A_{21} & 0 \\ 0 & \gamma_1^2 A_{12}^{\dagger} A_{12} \end{bmatrix} - \begin{bmatrix} I & 0 \\ 0 & I \end{bmatrix} < 0$$
  
$$\Leftrightarrow \bar{\sigma} \left( \begin{bmatrix} \gamma_2 I & 0 \\ 0 & \gamma_1 I \end{bmatrix} \begin{bmatrix} A_{21} & 0 \\ 0 & A_{12} \end{bmatrix} \right) < 1 \quad (21)$$

The stability condition (21) is stated formally below.

*Lemma 2:* Assume system (13) has been decomposed into the subsystems given in (14). Further assume that the subsystems are dissipative with respect to  $S_i = 0, Q_i = -I$ and  $R_i = \gamma_i^2 I$  where  $\gamma_i$  denotes the  $\mathcal{L}_2$ -norm of subsystem *i*. Then if max  $\{\gamma_2 \bar{\sigma}(A_{21}), \gamma_1 \bar{\sigma}(A_{12})\} < 1$  system (13) is asymptotically stable as verified by the Lyapunov function  $V_1(x_1) + V_2(x_2)$ .

Lemma 2 and Equation (21) provide a *nominal* stability test for the decomposed subsystems. What would be desirable it to determine the maximum  $\mathcal{L}_2$  gains (i.e.  $\gamma_i$ 's) such that (21) holds. Such a characterization would provide a robustness measure for the decomposed system. From the equivalence relation in (21) the maximum achievable  $\gamma$ 's denoted by  $\hat{\gamma}$  that satisfy the stability requirement in Lemma 2 are given by  $\hat{\gamma}_1 = \bar{\sigma}(A_{12})^{-1}$  and  $\hat{\gamma}_2 = \bar{\sigma}(A_{21})^{-1}$ .

2 are given by  $\hat{\gamma}_1 = \bar{\sigma}(A_{12})^{-1}$  and  $\hat{\gamma}_2 = \bar{\sigma}(A_{21})^{-1}$ . If we consider more generic supply rates with  $(-\kappa I, 0, \kappa \gamma^2 I)$ ,  $\kappa > 0$  instead of  $(-I, 0, \gamma^2 I)$  then it is possible to strengthen Lemma 2. Observe that if a system is dissipative w.r.t.  $(-I, 0, \gamma^2 I)$  then it is also dissipative w.r.t.  $(-\kappa I, 0, \kappa \gamma^2 I)$  for any  $\kappa > 0$ .

*Lemma 3:* If there exists a scalar  $\kappa > 0$  such that  $\Sigma_1$  and  $\Sigma_2$  are dissipative w.r.t.  $(-\kappa I, 0, \kappa \gamma_1^2 I)$  and  $(-\kappa I, 0, \kappa \gamma_2^2 I)$  respectively then system (13) is stable if max  $\{\kappa \gamma_1 \bar{\sigma}(A_{21}), \kappa^{-1} \gamma_2 \bar{\sigma}(A_{12})\} < 1$ .

Note that when  $\bar{\sigma}(A_{21}) = \bar{\sigma}(A_{12})$ , Lemma 3 is equivalent to the small gain condition.

#### C. Input Strong Passivity

A system of the form (5) is said to be *input strongly* passive if  $\int_0^T u^T y - \eta u^T u \, dt \ge 0$  for all  $T \ge 0$ , x(0) = 0 with a dissipation rate  $\eta < 0$  [10]. The maximum dissipation achievable is the largest  $\eta$  such that the integral inequality above holds, this can be computed by solving the LMI

$$\begin{array}{ll} \max & \eta \\ \text{s.t.} & \left[ \begin{array}{c} A^{\top}P + PA & PB - C^{\top} \\ B^{\top}P - C & 2\eta I \end{array} \right] < 0 \\ P > 0, \quad \eta < 0. \end{array}$$

Setting  $\tau = -2\eta$  we obtain a Q, S, R supply rate of  $(0, I, -\tau I)$ . Substituting the appropriate supply rate functions  $(0, I, -\tau_i I)$  corresponding to  $\Sigma_1$  and  $\Sigma_2$  into Equation (16) the requirement for  $V_1(x_1) + V_2(x_2)$  to be a composite Lyapunov function for (13) becomes

$$\begin{bmatrix} \tau_2 A_{21}^{\top} A_{21} & A_{12} + A_{21}^{\top} \\ A_{12}^{\top} + A_{21} & \tau_1 A_{12}^{\top} A_{12} \end{bmatrix} < 0$$
(23)

when  $H_{11} = H_{22} = 0$ . The LMI (23) is never feasible as the elements on the diagonal will always be positive definite. However if we allow for a feedback term in each subsystem by modifying the decomposition according to (18) then it is possible to make (16) negative definite. The modified LMI is easily obtained from (16) but is omitted for lack of space.

Although (23) can never be satisfied it does provide useful insight into the decomposition problem. Factoring (23) into

$$\begin{bmatrix} A_{21}^{\top} & 0\\ 0 & A_{12}^{\top} \end{bmatrix} \begin{bmatrix} \tau_2 & 0\\ 0 & \tau_1 \end{bmatrix} \begin{bmatrix} A_{21} & 0\\ 0 & A_{12} \end{bmatrix} + \begin{bmatrix} 0 & \Gamma\\ \star & 0 \end{bmatrix} < 0$$

where  $\Gamma = A_{12} + A_{21}^{\top}$ , we see that a heuristic that could be incorporated into a decomposition scheme that aims to provide a composite Lyapunov function for (13) with dissipative subsystems is

$$\min_{A_{21},A_{12}} \bar{\sigma} \left( \left[ \begin{array}{cc} \sqrt{\tau_2} & 0\\ 0 & \sqrt{\tau_1} \end{array} \right] \left[ \begin{array}{cc} A_{21} & 0\\ 0 & A_{12} \end{array} \right] \right),$$

where  $A_{12}, A_{21}$  are permuted blocks of A. The idea of synthesizing a system decomposition will be discussed further in the following sections.

#### **IV. DECOMPOSITION**

In this section we study the decomposition of interconnected dissipative subsystems. Consider a subsystem  $\Sigma_i$ described by (5) with a nonnegative storage function  $V_i(x_i)$ satisfying (4). Such storage is a scalar measure of the subsystem's state, which could be thought of as the amount of "abstract energy" stored by the subsystem in its internal state  $x_i$ .

Assuming  $\Sigma_i$  to be dissipative with respect to a given supply rate, a dissipation inequality indicates that the supply rate upper bounds the rate of change of the storage and thus, indirectly, the change of the subsystem's state. By estimating various forms of supply interchanges between subsystems we can evaluate which subsystems interact most strongly with each other using the supply rate upperbound as an indication of the worst case (strongest) interaction.

### A. Undirected Supply Measure

With the interconnection of subsystems  $\Sigma_1$  and  $\Sigma_2$  and when Assumption 1 holds, having the total supply rate given by  $w_1(u_1, y_1) + w_2(u_2, y_2) = -y^{\top} \hat{Q} y \approx 0$  implies that, over the trajectories of the global system, the interconnections H supply to the subsystems is small. Hence  $w_1(u_1, y_1) + w_2(u_2, y_2)$  could be seen as a measure of *undirected interaction*, indicating how relevant the interconnection is to the global system dynamics. Additionally, from Lemma 1 having  $w_1(\cdot) + w_2(\cdot) < 0$  implies stability of the global system.

*Remark 3:*  $\Sigma_1$  and  $\Sigma_2$  could be connected to several other subsystems. For the previous discussion to hold, one should constrain the supply rate to be separable along the different edges.

The above remark imposes the following constraint:

Assumption 2: Define  $\mathcal{E}_c \subset \mathcal{E}$  as the set of edges interconnecting K subsystems. For each subsystem  $\Sigma_i$  we assume the supply rate is separable along the edges, which implies  $w_i(u_i, y_i) = \sum_{e_k \in \mathcal{E}_c} w_{e_k}^i(u_i, y_i)$ .

*Remark 4:* Define  $\mathcal{E}_c^{ij} \subseteq \mathcal{E}_c$  as the set of edges connecting  $\Sigma_i$  and  $\Sigma_j$ . The measure for the

undirected interaction between  $\Sigma_i$  and  $\Sigma_j$  is then  $\sum_{e_k \in \mathcal{E}_c^{ij}} [w_{e_k}^i(u_i, y_i) + w_{e_k}^j(u_j, y_j)].$ 

Take the RC-network described previously. The metric discussed in this section corresponds to the electric energy dissipated in the interconnecting resistors for Observation 1 and to the electric energy dissipated on the interconnecting and boundary resistors for Observation 2.

#### B. Directed Supply Measure

*Remark 5:* A directed supply measure can also be derived using similar arguments as above, resulting in  $w_1(u_1, y_1)$ being a measure for *directed interaction* from  $\Sigma_2$  and the interconnection to  $\Sigma_1$ , see [?].

# NEED TO ADD REFERENCE TO THE TECHNICAL REPORT.

# C. Computing Edge Weights for Stability

We now discuss how the previously described measures of interaction, which are time varying functions of the system state, can be condensed to a representative static value for use in a decomposition algorithm such as the one in [5]. Ideally we would like to find "good" decompositions that satisfy the stability criteria defined in Section III.

Consider the global interconnected subsystem (7), with no assumption of stability. Define for each edge  $e_i$  the supply rate function

$$w_{e_i}(u, y) = y^{\top} Q_{e_i} y + 2u^{\top} S_{e_i} y + u^{\top} R_{e_i} u = y^{\top} \hat{Q}_{e_i} y$$
(24)

where u has been eliminated using the interconnection u = -Hy and  $\hat{Q}_{e_i}$  symmetric, corresponding to either the directed or the undirected interaction measure. For instance, taking the RC-network in Figure 1 and considering the undirected interaction measure for the supply rate from Observation 1,  $w_i(u_i, y_i) = u_i^\top y_i$ , we have

$$w_{e_1}(y_1, y_2) = \begin{bmatrix} y_1 \\ y_2 \end{bmatrix}^\top \begin{bmatrix} \frac{1}{\hat{R}_1} & -\frac{1}{\hat{R}_1} \\ -\frac{1}{\hat{R}_1} & \frac{1}{\hat{R}_1} \end{bmatrix} \begin{bmatrix} y_1 \\ y_2 \end{bmatrix} = \frac{(y_1 - y_2)^2}{\hat{R}_1},$$

which corresponds to the electric power dissipated by the resistor  $\widehat{R}_1$ .

As mentioned in Section III and IV-A, the undirected interaction measure is also related to stability. In fact given a cut  $\mathcal{E}_c$  and Assumption 2, if the subsystems  $\Sigma_1$  and  $\Sigma_2$  are dissipative with nonnegative storage functions characterized by  $P_1$  and  $P_2$  respectively, then we have

$$x^{\top} \left[ (A - BHC)^{\top} P + P(A - BHC) \right] x \leq y^{\top} \left( \sum_{e_k \in \mathcal{E}_c} \hat{Q}_{e_k} \right) y,$$
(25)

with  $P = \text{diag}(P_1, P_2)$ , from which it follows that the global system is stable if  $y^{\top}(\sum_{e_k \in \mathcal{E}_c} \hat{Q}_{e_k})y < 0$  (see Lemma 1).

As such a cut is not known *a priori*, we provide heuristics to compute appropriate edge weights. Let  $y_{e_k}$  be the output of the two nodes incident to the edge  $e_k$  and define  $\tilde{Q}_{e_k}$ such that  $y^{\top}\hat{Q}_{e_k}y = y_{e_k}^{\top}\tilde{Q}_{e_k}y_{e_k}$ . For a suitable permutation yielding  $\Pi y = [y_{e_k}^{\top}y_{\mathcal{E}/e_k}^{\top}]^{\top}$  we have  $\tilde{Q}_{e_k}$  as the first diagonal block of  $\Pi \hat{Q}_{e_k} \Pi^{-1}$ . For instance, in the RC-network in Figure 1 we have  $\tilde{Q}_{e_k} = \hat{Q}_{e_k}$ . A sufficient condition for  $\hat{Q}_{\mathcal{E}_c} = \sum_{e_k \in \mathcal{E}_c} \hat{Q}_{e_k} < 0 \text{ is to require } \tilde{Q}_{e_k} < 0 \quad \forall e_k \in \mathcal{E}.$  Furthermore, note that (25) may be thought as the inclusion of an ellipsoid,  $-x^{\top} \left[ (A - BHC)^{\top}P + P(A - BHC) \right] x$ , by another ellipsoid  $-x^{\top}C^{\top} (\sum_{e_k \in \mathcal{E}_c} \hat{Q}_{e_k}) Cx$ , where the latter corresponds to a sum of ellipsoids. Since the former ellipsoid is only known after the cut, one would like the latter ellipsoid to be as large as possible, as this would increase the set of matrices  $P_1$  and  $P_2$  for which such inclusion holds. Therefore, assuming  $\hat{Q}_{\mathcal{E}_c} < 0$  and denoting  $\mathcal{P}_{\mathcal{E}_c} = \{y: -y^{\top}\hat{Q}_{\mathcal{E}_c}y \leq 1\}$  as the ellipsoid associated with a given cut  $\mathcal{E}_c$ , a suitable partitioning algorithm would solve  $\max_{\mathcal{E}_c} \operatorname{vol}(\mathcal{P}_{\mathcal{E}_c})$  where  $J(\mathcal{E}_c) = \operatorname{vol}(\mathcal{P}_{\mathcal{E}_c})$  is the utility of a cut  $\mathcal{E}_c$ . Combining these two features, and denoting  $\mathcal{P}_{e_k}$  as the ellipsoid defined by  $\tilde{Q}_{e_k}$ , the edge weights  $J(e_k) = \operatorname{vol}(\mathcal{P}_{e_k}^e)$  may be computed by solving

$$\max_{Q_{e_k} < 0, R_{e_k} = R_{e_k}^{\top}, S_{e_k}} \quad \text{vol}(\mathcal{P}_{e_k})$$

which is not a convex problem in  $Q_{e_k}$ ,  $S_{e_k}$ , and  $R_{e_k}$ . The partitioning algorithm would then choose a set of edges forming a cut  $\mathcal{E}_c$  such that  $\sum_{e_k \in \mathcal{E}_c} J(e_k)$  is maximized. <u>The volume</u> of an ellipsoid  $\mathcal{P}_{e_k}$  is proportional to

The volume of an ellipsoid  $\mathcal{P}_{e_k}$  is proportional to  $\sqrt{\det(-\tilde{Q}_{e_k}^{-1})}$ , hence the previous problem will yield a solution such that  $\det(-\tilde{Q}_{e_k})$  is minimized, which implies  $y_{e_k}^{\top}\tilde{Q}_{e_k}y_{e_k}\approx 0$ . Therefore, since  $\det(-\tilde{Q}_{e_k})$  is the product of the eigenvalues of  $-\tilde{Q}_{e_k}$ , we can instead consider the convex problem

$$\min_{\substack{Q_{e_k} < 0, R_{e_k} = R_{e_k}^{\top}, S_{e_k}}} \qquad \lambda_{max}(-\tilde{Q}_{e_k})$$

which is related to maximizing the diameter of  $\mathcal{P}_{e_k}$ , and take  $J(e_k) = 1/\lambda_{max}(-\tilde{Q}_{e_k}^*)$ . Note that this also relates to finding weakly interacting subsystems based on the undirected measure. Therefore, large values of  $J(e_k)$  indicate that this is a good edge to cut in a decomposition algorithm and will also help in verifying stability using the criteria in Section III.

#### D. Computing Edge Weights for Weakly Connected Systems

Assume now the global system is stable and we want to decompose it into subsystems that interact weakly over time, for example to facilitate the design of distributed controllers. Consider the system (7) and define for each edge the supply rate function (24). Since there are no external inputs,  $w_{e_i}$  is a function of the initial condition  $x_0$  and time. Hence one needs to evaluate these functions to compute a static value measuring the interactions over time. Instead the *total supply* defined as  $W_{e_i}(x_0) \triangleq \int_0^\infty w_{e_i}(t) dt$  is used and the edge weight is computed by evaluating  $W_{e_i}(x_0)$  for the relevant initial conditions. The following result allows us to compute  $W_{e_i}(x_0)$  for a given initial condition:

Proposition 1: Assuming the global system (7) is stable, for a given initial condition  $x_0$  we have  $W_{e_i}(x_0) = x_0^{\top} T_{e_i} x_0$ , where  $T_{e_i}$  is the Gramian matrix satisfying the Lyapunov equation  $(A - BHC)^{\top} T_{e_i} + T_{e_i} (A - BHC) + C^{\top} \hat{Q}_{e_i} C = 0$ .

equation  $(A - BHC)^{\top}T_{e_i} + T_{e_i}(A - BHC) + C^{\top}\hat{Q}_{e_i}C = 0.$ *Proof:* We have  $W_{e_i}(x_0) = \int_0^\infty y(t)^{\top}\hat{Q}_{e_i}y(t) dt = x_0^{\top}Tx_0$ , where  $\bar{A} = A - BHC$  and  $T = \int_0^\infty e^{(\bar{A}^{\top}t)}C^{\top}\hat{Q}_{e_i}Ce^{(\bar{A}t)} dt.$  Note that the expression for T resembles the well-known observability Gramian. The rest of the proof follows the characterization of the observability Gramian found in [14]. Note that for finite time horizons  $W_{e_i}(x_0)$  can be computed

by means of simulation as an alternative to solving the Gramian.

Recalling the decomposition's objective, a partitioning algorithm would select a set of edges  $\mathcal{E}_c$  forming a cut such that  $\sum_{e_i \in \mathcal{E}_c} W_{e_i}(x_0)$  is close to zero, thus solving  $\min_{\mathcal{E}_c} |\sum_{e_i \in \mathcal{E}_c} W_{e_i}(x_0)|$  where we define  $J(\mathcal{E}_c) = |\sum_{e_i \in \mathcal{E}_c} W_{e_i}(x_0)|$  as the cost of a cut for a given initial condition. We now analyze the evaluation of  $W_{e_i}(x_0)$  for two different sets of initial conditions.

1) Worst-case initial condition: For a given cut  $\mathcal{E}_c$ , the worst-case initial condition is the one maximizing  $|\sum_{e_i \in \mathcal{E}_c} W_{e_i}(x_0)|$  and the cut cost would be given by  $J(\mathcal{E}_c) = \max_{||x_0||=1} |\sum_{e_i \in \mathcal{E}_c} W_{e_i}(x_0)|$ . Since we do not know the set  $\mathcal{E}_c$  a priori, the edge weights are also unknown, which would require a combinatorial approach to solve this partitioning algorithm. A possible relaxation decoupling the edge weights from the cut can be made based on the following inequality  $\max_{x_0} |\sum_{e_i \in \mathcal{E}_c} W_{e_i}(x_0)| \leq \sum_{e_i \in \mathcal{E}_c} \max_{x_{0_i}} |W_{e_i}(x_{0_i})|$ . Note that in the right hand side the initial condition  $x_{0_i}$  is dependent only the edge  $e_i$ . Therefore, by defining the new cost function  $\overline{J}(\mathcal{E}_c) = \sum_{e_i \in \mathcal{E}_c} \max_{x_{0_i}} |W_{e_i}(x_{0_i})|$ , we obtain weights that only depend on each particular edge,  $\overline{J}(e_i) = \max_{x_{0_i}} |W_{e_i}(x_{0_i})|$ , and an upper bound on the edge cost  $J(\mathcal{E}_c)$ . The weight  $\overline{J}(e_i)$  can be computed by solving  $\max_{||x_0||=1} |x_0^\top T_{e_i}x_0|$  where  $T_{e_i}$  is given by Proposition 1.

Remark 6: Since  $C^{\top}\hat{Q}_{e_i}C$  is symmetric, T is also symmetric and thus we have  $\max_{x_0} |x_0^{\top}Tx_0| = \max_i |\lambda_i(T)| = |\lambda^*(T)|$ , where  $\{\lambda_i(T)\}$  are the eigenvalues of T. Computing the eigenvalue value decomposition of T, we conclude  $x_0^*$  is given by the eigenvector associated with  $\lambda^*(T)$ .

2) Gaussian initial condition: We now consider a stochastic description of the initial condition for the global system. Let  $x_0 \sim \mathcal{N}(\bar{x}, \Omega)$ . From Proposition 1 it follows that the total supply  $W_{e_i}(x_0) = x_0^{\top} T_{e_i} x_0$  is a random variable. Hence the cost of a given cut  $\mathcal{E}_c$  is  $J(\mathcal{E}_c) = |\mathbb{E}_{x_0}[\sum_{e_i \in \mathcal{E}_c} W_{e_i}(x_0)]|$ . Using the triangle inequality we obtain the following upper bound of the cut cost  $J(\mathcal{E}_c) \leq \sum_{e_i \in \mathcal{E}_c} |\mathbb{E}_{x_0}[W_{e_i}(x_0)]| = \tilde{J}(\mathcal{E}_c)$ . Hence we assign  $\tilde{J}(e_i) = |\mathbb{E}_{x_0}[x_0^{\top} T_{e_i} x_0]|$  as the weight for  $e_i$ , which may be computed using the following result:

Proposition 2: Given  $x_0 \sim \mathcal{N}(\bar{x}, \Omega)$  we have  $\mathbb{E}_{x_0}[x_0^\top T_{e_i} x_0] = \bar{x}^\top T_{e_i} \bar{x} + \operatorname{trace}(T_{e_i} \Omega).$ 

# E. Decomposition Methods

Given the aforementioned methods to compute static edge weights, a system decomposition algorithm based on the directed and undirected interaction measures is described.

The directed interaction measure provides two weights, one for each edge direction. Hence it is suitable for clustering algorithms where a given set of nodes  $\mathcal{V}_0$  is of interest and we want to find  $\mathcal{V}_i$  such that  $\mathcal{V} = \mathcal{V}_i \cup \mathcal{V}_j, \mathcal{V}_0 \subseteq \mathcal{V}_i$ , and  $\Sigma_i$  is not affected much by  $\Sigma_j$ . A possible algorithm to accomplish this tasks proceeds as follows:

- Set V<sub>i</sub> = V<sub>0</sub> and define E<sub>c</sub> as the edge set connecting nodes from V<sub>i</sub> to V<sub>i</sub>;
- Compute the directed weight from V<sub>j</sub> to V<sub>i</sub> for each edge e<sub>k</sub> ∈ E<sub>c</sub>;
- 3) Pick the set of nodes from V<sub>j</sub> that have the largest directed weight, V
  <sub>j</sub>, and set V<sup>+</sup><sub>i</sub> = V<sub>i</sub> + V
  <sub>j</sub>;
   4) Set V<sub>i</sub> = V<sup>+</sup><sub>i</sub>, define the new cut set E<sub>c</sub>, and repeat
- 4) Set  $\mathcal{V}_i = \mathcal{V}_i^+$ , define the new cut set  $\mathcal{E}_c$ , and repeat from 2 until the interaction measure is below the tolerance level.

The undirected interaction measure provides a single weight,  $W_{e_k}$  which can be readily incorporated into the framework presented in [5].

#### V. EXAMPLE

Consider an RC-network described by the graph in Figure 1 with dynamics given by (11). Let each node have unit capacitance and resistance and let  $[W]_{ii} = 1/\hat{R}_i = 0.1 \forall e_i \in \mathcal{E}_c^{12}$  and  $[W]_{ii} = 1 \forall e_i \notin \mathcal{E}_c^{12}$ . For each node  $v_i$ , consider the supply rate defined in Observation 1,  $w_i(u_i, y_i) = u_i^\top y_i$ . Recalling that  $u_i$  is the input current to node i and  $y_i = x_i$  is the voltage at the corresponding capacitor, from Kirchhoff's Current Law we conclude that the supply rate is separable along the edges connected to  $v_i$ , since  $w_i(u_i, y_i)$  is the sum of the input power from each edge. Hence Assumption 2 holds. Following the steps in Section IV.E for the *undirected* measure and the *worst-case* initial condition approach we compute the edge weights, which correspond to the electric power dissipated at each edge resistor. For the dashed edges we obtain the weights

 $\bar{J}(e_i \in \mathcal{E}_c^{12}) = [0.0579, 0.0625, 0.0693, 0.0623]^{\top}$ , while  $\min_i \bar{J}(e_i \notin \mathcal{E}_c^{12}) = 0.4016$ . Applying a spectral graph decomposition algorithm with these edge weights we obtain  $\mathcal{E}_c^{12}$  as the cut set, as shown in Figure 1.

Considering instead the undirected measure with Gaussian initial condition  $x_0 \sim \mathcal{N}(0, I)$ , we obtain the following weights  $\tilde{J}(e_i \in \mathcal{E}_c^{12}) = [0.0636, 0.0670, 0.0727, 0.0671]^\top$ , and  $\min_i \tilde{J}(e_i \notin \mathcal{E}_c^{12}) = 0.4278$ . As before, the cut set obtained after spectral decomposition is  $\mathcal{E}_c^{12}$ .

# VI. CONCLUSIONS

It has been shown how the supply rates of dissipative dynamical systems can be used as a metric for measuring subsystem interaction strength in a networked system. Furthermore, based upon this metric an algorithm for decomposing a networked system was presented and illustrated on a 20 node RC circuit. We also provided stability criteria for the decomposed system based on passivity and bounded gain.

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